

Derived modular operads and metric ribbon graphs

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- 1 Motivation
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- 3 Grafting flags and contracting edges
- 4 Cyclic operads and modular operads
- 5 W -construction and moduli spaces

Definition (moduli space for oriented surfaces/ribbon graphs)

- $\mathcal{M}_{g,n}$ moduli space of *hyperbolic metrics* on a closed oriented surface of *genus g* with *n punctures*.
- $\mathcal{MRG}_{g,n}$ moduli space of *admissible metrics* on a ribbon graph of type (g, n) .

Theorem (Mumford, Harer, Penner, Strebel, Kontsevich)

$\mathcal{M}_{g,n} \simeq \mathcal{MRG}_{g,n}$ if $n > 0$ and $2 - 2g < n$.

Purpose of the talk (cf. Igusa, Costello, Giansiracusa)

Realise $\mathcal{MRG}_{g,n}$ as a “*derived*” *modular envelope* of the cyclic operad CycAss of planar structures.

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Definition (graph)

A *graph* is quadruple (V, F, ∂, ι) consisting of a vertex-set V , a flag-set F , a boundary map $\partial : F \rightarrow V$ and an involution $\iota : F \rightarrow F$

A two-element orbit of ι is called an *edge*, any fixpoint of ι an *outer flag*. For each vertex $v \in V$ the set $\partial^{-1}(v)$ is the flag-set of the *star* $*_v$ at v . Each graph decomposes into stars.

Definition (ribbon graph)

- A *ribbon graph* is a graph equipped with a cyclic flag-ordering for each vertex-star, i.e. equipped with a permutation $\sigma_0 : F \rightarrow F$ whose orbits correspond to the vertex-stars.
- A *boundary cycle* of a ribbon graph is an orbit of $\sigma_\infty = \sigma_0 \iota$.
- A ribbon graph is said to be of type (g, n) if it has n boundary cycles and $2 - 2g = \#V - \#E + n$.

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Proposition (unflagged ribbon graphs)

For a finite graph G without outer flags TFAE

- G comes equipped with a ribbon structure of type (g, n) ;
- $|G|$ embeds into a closed oriented surface S of genus g such that $S - |G|$ is a disjoint union of n discs.

Corollary (planar graphs)

A graph is planar (resp. a planar tree) if and only if it carries a ribbon structure of type $(0, n)$ (resp. $(0, 1)$).

Remark (flagged ribbon graphs - doubling construction)

- The outer flags of a ribbon graph G form a polycyclic set.
- Gluing $G \cup G^{op}$ along outer flags yields an *unflagged* ribbon graph with an "orientation-reversing" involution.
- This yields an equivalence between *flagged* ribbon graphs and certain *involutive* unflagged ribbon graphs.

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Definition (Borisov-Manin graph morphism)

A morphism $G = (V, F, \partial, \iota) \rightarrow G' = (V', F', \partial', \iota')$ consists of a triple $(\phi_V, \phi^F, \gamma_\phi)$ where

- $\phi_V : V \rightarrow V'$ surjection;
- $\phi^F : F' \rightarrow F$ outer flag preserving injection;
- γ_ϕ fixpoint-free involution on $F_{out} \setminus \phi^F(F'_{out})$

such that $\partial' = \phi_V \partial \phi^F$ and $\phi^F \iota' = \iota \phi^F$.

A morphism is called a *grafting* if ϕ_V and ϕ^F are bijections. A morphism is called a *virtual contraction* if there is no grafting.

Lemma (unique factorisation system)

Each Borisov-Manin graph morphism factors essentially uniquely as a grafting followed by a virtual contraction.

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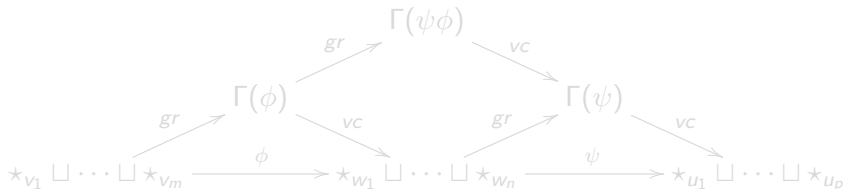
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Definition (Feynman category \mathfrak{F}_{agg} of aggregates)

The objects of \mathfrak{F}_{agg} are finite coproducts of corollas.

The morphisms of \mathfrak{F}_{agg} are Borisov-Manin graph morphisms.



where $\Gamma(\psi\phi)$ is the graph obtained by *inserting* the n formal components of $\Gamma(\phi)$ into the n vertices of $\Gamma(\psi)$.

Remark (insertional classes \mathcal{C} of graphs vs Feynman subcategories)

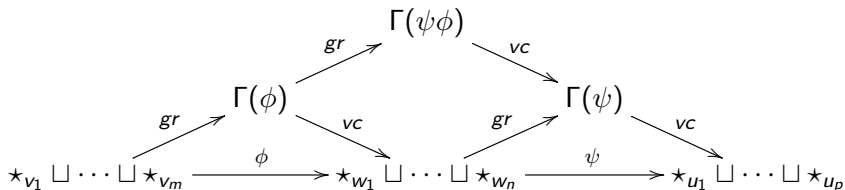
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For $\mathcal{C} = (\text{connected graphs})$, we get a Feynman category \mathfrak{F}_{ctd} where the formal components of $\Gamma(\phi)$ coincide with its path components!

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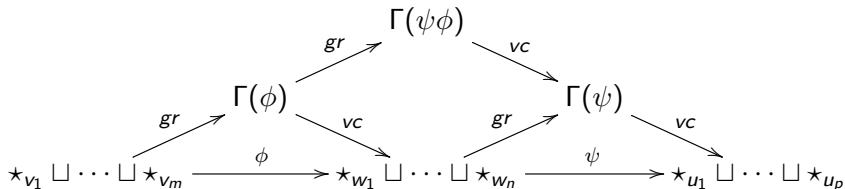
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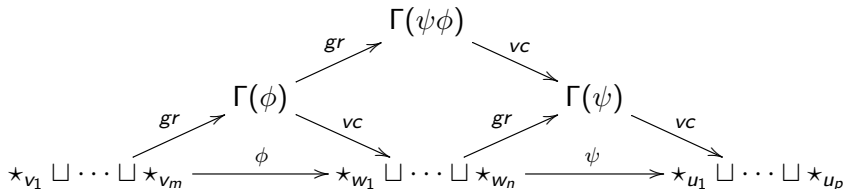
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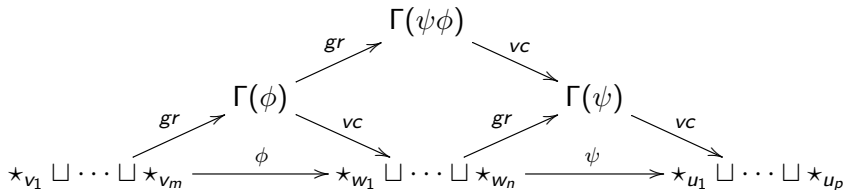
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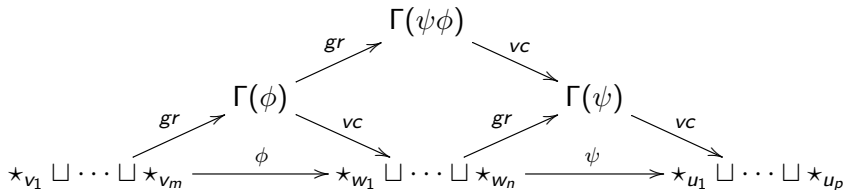
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where $\Gamma(\psi\phi)$ is the graph obtained by *inserting* the n formal components of $\Gamma(\phi)$ into the n vertices of $\Gamma(\psi)$.

Remark (insertional classes \mathfrak{C} of graphs vs Feynman subcategories)

Each \mathfrak{C} induces a Feynman subcategory $\mathfrak{F}_{\mathfrak{C}} \subset \mathfrak{F}_{agg}$ and vice-versa.

For $\mathfrak{C} = (\text{connected graphs})$, we get a Feynman category \mathfrak{F}_{ctd} where the formal components of $\Gamma(\phi)$ coincide with its path components !

Proposition (Street-Walters, K-B)

Each (Feynman) functor factors essentially uniquely as a connected (Feynman) functor followed by a covering.

Definition

Define $\mathfrak{F}_{cyc} = \mathfrak{F}_{(trees)}$ and \mathfrak{F}_{mod} by the connected/covering factorisation $\mathfrak{F}_{(trees)} \xrightarrow{\text{connected}} \mathfrak{F}_{mod} \xrightarrow{\text{covering}} \mathfrak{F}_{(ctd\ graphs)}$.

Proposition (Getzler-Kapranov)

Symmetric monoidal functors out of \mathfrak{F}_{cyc} (resp. \mathfrak{F}_{mod}) are non-unital cyclic (resp. modular) operads.

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For each $P \in \text{Func}_{\otimes}(\mathfrak{F}, \text{Sets})$ there is a covering $\mathfrak{F}_{dec(P)} \rightarrow \mathfrak{F}$ inducing an equivalence $\text{Func}_{\otimes}(\mathfrak{F}, \text{Sets})/P \simeq \text{Func}_{\otimes}(\mathfrak{F}_{dec(P)}, \text{Sets})$.

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Example (Computation of $j_! \text{CycAss}$, cf. Markl, Doubek, B-K)

$$(j_! \text{CycAss})(\star_{\gamma, \nu})$$

$$= \text{colim}_{j(-) \downarrow \star_{\gamma, \nu}} \text{CycAss}(-)$$

= (equ. cl. of ribbon graphs with γ loops and ν outer flags)

= (topological types (g, n, S_1, \dots, S_k) of bordered oriented surf.)

Corollary (B-K)

The morphisms of the Feynman category $\mathfrak{F}_{\text{surf-mod}}$ are genus/puncture-labelled *polycyclic* graphs.

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A Feynman category \mathfrak{F} is *cubical* if there is a degree function $\deg : \text{Mor}(\mathfrak{F}) \rightarrow \mathbb{N}_0$ such that

- $\deg(\phi \otimes \psi) = \deg(\phi) + \deg(\psi)$
- $\deg(\phi \circ \psi) = \deg(\phi) + \deg(\psi)$
- Degree 0 morphisms are invertible
- Each degree n morphism factors (up to iso) in $n!$ ways into degree 1 morphisms “compatibly with composition”

Remark

The Feynman categories \mathfrak{F}_{ctd} , \mathfrak{F}_{cyc} , \mathfrak{F}_{mod} , $\mathfrak{F}_{plan-cyc}$, $\mathfrak{F}_{surf-mod}$ are all cubical. The degree of ϕ is the number of edges of the associated graph $\Gamma(\phi)$. However, \mathfrak{F}_{agg} is *not* cubical.

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Definition (W -construction for $P \in \text{Func}_{\otimes}(\mathfrak{F}, \text{sSets})$, \mathfrak{F} cubical)

$$(W_{\mathfrak{F}}P)(B) = \Delta[1]^{\text{deg}(-)} \otimes_{\mathfrak{F} \downarrow B} P \circ \text{dom}(-)$$

where $\Delta[1]^{\text{deg}(-)}$ are cube embeddings via 0-face.

Proposition (K-Ward, cf. Boardman-Vogt, B-Moerdijk for $\mathfrak{F}_{\text{sym}}$)

For any cubical Feynman category \mathfrak{F} , the category $\text{Func}_{\otimes}(\mathfrak{F}, \text{sSets})$ admits a *transferred model structure*. If P has an underlying cofibrant \mathcal{V} -collection then $W_{\mathfrak{F}}P$ is *cofibrant* in $\text{Func}_{\otimes}(\mathfrak{F}, \text{sSets})$.

Lemma (relative W -construction for cubical $f : \mathfrak{F} \rightarrow \mathfrak{F}'$)

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Example (cubically subdivided convex polytopes)

- $W_{\tilde{\mathfrak{S}}_{\text{sym}}}(\text{Ass})(\text{rooted corolla}) = \text{associahedron}$
- $W_{\tilde{\mathfrak{S}}_{\text{cyc}}}(\text{CycAss})(\text{corolla}) = \text{cyclohedron}$

Theorem (B-K)

$$|J_!(W_{\tilde{\mathfrak{S}}_{\text{plan-cyc}}} \mathbf{1})(g, n, S_{\bullet})| \simeq \text{MRG}_{g,n,S_{\bullet}}$$

Definition (Igusa)

The category $\text{rb}_{g,n,S_{\bullet}}$ has at least trivalent ribbon graphs of type (g, n, S_{\bullet}) as objects, and ribbon contractions as morphisms.

Theorem (Igusa, B-K)

$$\text{MRG}_{g,n,S_{\bullet}} \simeq |\text{nerve}(\text{rb}_{g,n,S_{\bullet}})|$$

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$$|J_!(W_{\mathfrak{S}_{plan-cyc}} \mathbf{1})(g, n, S_{\bullet})| \simeq \text{MRG}_{g,n,S_{\bullet}}$$

Definition (Igusa)

The category $\text{rb}_{g,n,S_{\bullet}}$ has at least trivalent ribbon graphs of type (g, n, S_{\bullet}) as objects, and ribbon contractions as morphisms.

Theorem (Igusa, B-K)

$$\text{MRG}_{g,n,S_{\bullet}} \simeq |\text{nerve}(\text{rb}_{g,n,S_{\bullet}})|$$

Thank you !